

Completely Mixed Bi-matrix Games Without Restrictions on Payoffs at Equilibrium

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Abstract

We consider bi-matrix games of the type discussed in Raghavan (1970) and Oviedo (1996). A strategy profile is said to be completely mixed if it assigns positive probabilities to all pure strategies for both players and a bi-matrix game is said to be completely mixed if all equilibria of the game are completely mixed. Our first theorem in this note extends necessary conditions for completely mixed bi-matrix games that comprise theorem 1 in Raghavan (1970) and prove that theorem 1 in Raghavan (1970) holds without assuming zero-valued equilibrium pay-offs. Results, related to the ranks of the matrices of a bi-matrix game are not “obviously independent” of the kind of restrictions on payoffs invoked in Raghavan (1970) and Oviedo (1996). An immediate corollary of our first theorem is that for completely mixed two-person zero-sum (TPZS) games with value zero, the pure-strategy pay-off matrices are square matrices with the rank of the matrices being one less than the dimension of the matrices. We apply this corollary to obtain a complete characterization of all completely mixed TPZS games that have value zero. This characterization is different from the characterization available in Kaplansky (1945). The Complementary Slackness Condition for bi-matrix games play a very useful and important role in our analysis.

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1. Notation: For any positive integer ℓ let $\Delta^{\ell-1} = \{x \in \mathbb{R}_+^\ell \mid \sum_{j=1}^{\ell} x_j = 1\}$ and for $j \in \{1, \dots, \ell\}$, let $E^{(\ell,j)}$ be the ℓ -dimensional column vector whose j^{th} coordinate is 1 and all other coordinates are equal to 0. Let $E^{(\ell)} = \sum_{j=1}^{\ell} E^{(\ell,j)}$. $E^{(\ell)}$ is said to be the ℓ -dimensional sum vector.

For positive integers m, n , and any $m \times n$ matrix P , for all $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, the $(i, j)^{\text{th}}$ term (term at the intersection of the i^{th} row and j^{th} column) of P is denoted by p_{ij} .

For $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, and any $m \times n$ matrix P , let P_i denote the i^{th} row of P and let P^j denote the j^{th} column of P .

For any $m \times n$ matrix P we will use P^T to denote its transpose.

2. Bi-matrix Games: For positive integers m, n , let A, B be two $m \times n$ matrices.

A **bi-matrix game** (interactive decision-making problem) with pay-off matrices A, B is a contest between two players- the **row player** and the **column player**- such that if for $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, the row player chooses row i and the column player chooses column j , then the pay-off to the row player is a_{ij} and the pay-off to the column player is b_{ij} .

A bi-matrix game with pay-off matrices A, B is denoted by the ordered pair (A, B) .

A bi-matrix game (A, B) is said to be a **two-person zero sum (TPZS) game** if $B = -A$.

It is customary to represent a TPZS game $(A, -A)$ by the matrix A , and just that.

A **strategy profile** is an ordered pair $(x, y) \in \Delta^{m-1} \times \Delta^{n-1}$, where x is the vector of probabilities with which the row-player chooses each strategy and y is the vector of probabilities with which the column-player chooses each strategy.

A very “witty” justification for randomized strategies is available in (page 17, in chapter 3 of) Washburn (2014): “It is the key to resolving the ‘... if he thinks that I think that he thinks that...’ impasse- you can’t outwit someone who refuses to think”.

An **equilibrium** for a bi-matrix game (A, B) is a strategy profile (x^*, y^*) that satisfies the following condition: For all $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, $A_i y^* \leq x^{*T} A y^*$ and $x^{*T} B_j \leq x^{*T} B y^*$.

Hence, (x^*, y^*) is an equilibrium if and only if $(x^*, y^*) \in \Delta^{m-1} \times \Delta^{n-1}$ and for all $(x, y) \in \Delta^{m-1} \times \Delta^{n-1}$ it is the case that $x^T A y^* \leq x^{*T} A y^*$ and $x^{*T} B y \leq x^{*T} B y^*$.

By the “Equivalence Theorem” of Mangasarian and Stone (1964), the set of equilibria of an $m \times n$ bi-matrix game is the projection into the $m + n$ dimensional Euclidean space of the set of solutions of a related quadratic programming problem.

Note 1: If C and D are $m \times n$ matrix such that for some real numbers, $c, d, c_{ij} = c$ and $d_{ij} = d$ for all $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, then it follows immediately that (x^*, y^*) is an equilibrium for (A, B) if and only if (x^*, y^*) is an equilibrium for $(A + C, B + D)$.

For the $m \times n$ bi-matrix game (A, B) and $(x, y) \in \Delta^{m-1} \times \Delta^{n-1}$ let $V(\text{row}(A, B), (x, y)) = x^T A y$ and $V(\text{column}(A, B), (x, y)) = x^T B y$.

Note 2: An immediate consequence of the definition of equilibrium is the following “**Complementary Slackness Condition for Bi-matrix Games**”: Suppose (x^*, y^*) is an equilibrium for the bi-matrix game (A, B) . (i) If for any $i \in \{1, \dots, m\}$ it is the case that $A_i y^* < x^{*T} A y^*$, then it must be the case that $x_i^* = 0$; and (ii) if for any $j \in \{1, \dots, n\}$ it is the case that $x^{*T} B_j < x^{*T} B y^*$, then it must be the case that $y_j^* = 0$.

It is well-known that associated with any TPZS game A , there is a real number $V(A)$, called the **value** of A , such if (x, y) is an equilibrium of A , then $x^T A y = V(A)$.

It has been known since the seminal work of John von-Neumann and Oscar Morgenstern entitled “Theory of Games and Economic Behaviour” that every TPZS game has at least one equilibrium. A “complete characterization theorem” for the set of all equilibria in terms of the solution sets of a linear programming and its dual, for a general class of bi-matrix games that includes the set all TPZS games as well as a large class of “prisoners’ dilemma” type games as a proper subset, is available as proposition 1 in Lahiri (2025). A very detailed treatment of

two-person zero-sum games is available in Washburn (2014). The scope of Washburn (2014) goes way beyond the two-person zero-sum (matrix) games, that we are concerned with here.

3. Linear Systems Associated with Bi-matrix games: We now provide a short lemma for bi-matrix games. Part (iii) of the lemma will play a significant role in proving uniqueness of equilibrium for completely mixed bi-matrix games.

Lemma 1: (i) If (x^*, y^*) is an equilibrium for the bi-matrix game (A, B) game then $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$ and there exists real numbers u^*, v^* such that $Ay^* \leq u^*E^{(m)}$, $B^T x^* \leq v^*E^{(n)}$. Further, $u^* = x^{*T}Ay^*$ and $v^* = x^{*T}By^*$.

(ii) If there exists $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$ and a real number u^* such that $Ay^* \leq u^*E^{(m)}$, $B^T x^* \leq -u^*E^{(n)}$ and $x^{*T}Ay^* = -x^{*T}By^*$ then (x^*, y^*) is an equilibrium for the bi-matrix game (A, B) .

(iii) If there exists $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$ and real numbers u^*, v^* such that $Ay^* = u^*E^{(m)}$, $B^T x^* = v^*E^{(n)}$, then (x^*, y^*) is an equilibrium for the bi-matrix game (A, B) .

Proof: (i) If (x^*, y^*) is an equilibrium for the bi-matrix game (A, B) , then with $u^* = x^{*T}Ay^*$ and $v^* = x^{*T}By^*$ we know that for all $i \in \{1, \dots, m\}$ and $j \in \{1, \dots, n\}$, $A_{ij}y^*_j \leq x^{*T}Ay^*$ and $x^{*T}B^j \leq x^{*T}By^*$, so that $Ay^* \leq u^*E^{(m)}$ and $B^T x^* \leq v^*E^{(n)}$. Further, $x^* \in \Delta^{m-1}$ and $y^* \in \Delta^{n-1}$.

(ii) Suppose that $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$ and there exists a real numbers u^* such that $Ay^* \leq u^*E^{(m)}$, $B^T x^* \leq -u^*E^{(n)}$ and $x^{*T}Ay^* = -x^{*T}By^*$.

Pre-multiplying the first inequality by x^{*T} and the second by y^{*T} we get $x^{*T}Ay^* \leq u^*$ and $x^{*T}By^* = y^{*T}B^T x^* \leq -u^*$.

Since, $x^{*T}By^* = -x^{*T}Ay^*$, we get $u^* \leq x^{*T}Ay^* \leq u^*$, so that $x^{*T}Ay^* = u^*$.

Thus, $x^{*T}By^* = -u^*$.

Since $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$, $Ay^* \leq u^*E^{(m)} = (x^{*T}Ay^*)E^{(m)}$ and $x^{*T}B \leq -u^*E^{(n)T} = (x^{*T}By^*)E^{(n)}$.

Thus, (x^*, y^*) is an equilibrium for (A, B) .

(iii) Now suppose $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$ and there exists real numbers u^*, v^* such that $Ay^* = u^*E^{(m)}$, $B^T x^* = v^*E^{(n)}$.

Thus, for any $x \in \Delta^{m-1}$ and $y \in \Delta^{n-1}$, $x^T Ay^* = u^* = x^{*T}Ay^*$ and $x^{*T}By = v^* = x^{*T}By^*$.

Thus, (x^*, y^*) is an equilibrium. Q.E.D.

4. Completely Mixed Equilibrium and Completely Mixed Bi-matrix Games: We now consider some properties for equilibrium that assigns positive probabilities to all pure strategies of both players and bi-matrix games all equilibria for which satisfy such a property.

An “equilibrium” (x^*, y^*) for a bi-matrix game (A, B) is said to be **completely mixed** if for all $i \in \{1, \dots, m\}$, $x_i^* > 0$ and for all $j \in \{1, \dots, n\}$, $y_j^* > 0$.

An immediate consequence of note 2 in the previous section is the following lemma.

Lemma 2: If (x^*, y^*) is a completely mixed equilibrium for the bi-matrix game (A, B) , then it must be the case that $Ay^* = (x^{*T}Ay^*)E^{(m)}$ and $x^{*T}B = (x^{*T}AB y^*)E^{(n)T}$.

A bi-matrix game (A, B) is said to be **completely mixed** if every equilibrium for (A, B) is completely mixed.

While the concept of completely mixed TPZS games originated in the work of Kaplansky (1945), the most recent significant results for such TPZS games as well as for completely mixed bi-matrix games can be found in Parthasarathy, Sharma and Sricharan (2020) and more recently in Parthasarathy, Gomatam and Kumar (2024). Some results in this note may overlap with results available in Oviedo (1996). We prove that if a bi-matrix game is completely mixed then it has a unique equilibrium. Unlike theorem 1 in Oviedo (1996) we do not assume that the pay-offs to the players in equilibrium is zero. We also extend several important conclusions in Raghavan (1970) and Oviedo (1996) to the set of all completely mixed bi-matrix games, without imposing any restrictions on the equilibrium pay-offs. Results, related to the ranks of the matrices of a bi-matrix game are not “obviously independent” of the kind of restrictions on payoffs invoked in Raghavan (1970) and Oviedo (1996).

Proposition 1: Suppose the bi-matrix game (A, B) is completely mixed. Then it has a unique equilibrium.

Proof: Suppose the bi-matrix game (A, B) is completely mixed. Thus, $A \neq 0 \neq B$.

Let (x^*, y^*) be a completely mixed equilibrium for (A, B) . By note 1, we may without loss of generality assume that $x^{*\top}Ay^* = 0 = x^{*\top}By^*$.

Since (x^*, y^*) is completely mixed, it must be the case that $x_i^* > 0$ for all $i \in \{1, \dots, m\}$ and $y_j^* > 0$ for all $j \in \{1, \dots, n\}$.

By lemma 2, it follows that $Ay^* = 0$ and $B^T x^* = 0$.

Thus, the columns of A are linearly dependent and the rows of B are linearly dependent.

Towards a contradiction suppose (x^0, y^0) is any other completely mixed equilibrium for (A, B) .

Thus, $(x^0, y^0) \neq (x^*, y^*)$, $x^0 \in \Delta^{m-1}$, $y^0 \in \Delta^{n-1}$, $x_i^0 > 0$ for all $i \in \{1, \dots, m\}$ and $y_j^0 > 0$ for all $j \in \{1, \dots, n\}$, and by lemma 2, $Ay^0 = (x^{0\top}Ay^0)E^{(m)}$ and $x^{0\top}B = (x^{0\top}By^0)E^{(n)\top}$.

Let $u^0 = x^{0\top}Ay^0$ and $v^0 = x^{0\top}By^0$.

Suppose $x^0 \neq x^*$. Thus, $\{i | x_i^0 > x_i^*\} \neq \emptyset$ and $\{i | x_i^0 < x_i^*\}$. Thus, $t = \max\{\frac{x_i^0}{x_i^*} | i \in \{1, \dots, m\}\} > 1 > \min\{\frac{x_i^0}{x_i^*} | i \in \{1, \dots, m\}\} = s > 0$.

Let $x' = x^0 - sx^*$. Clearly, $B^T x' = B^T x^0 - sB^T x^* = B^T x^0 = v^0 E^{(n)}$, $x' \in \mathbb{R}_+^m$.

Note that for all $i \in \{1, \dots, m\}$ satisfying $\frac{x_i^0}{x_i^*} = t$, we have $x_i^0 - sx_i^* = x_i^*(t - s) > 0$ and for all $i \in \{1, \dots, m\}$ satisfying $\frac{x_i^0}{x_i^*} = s$, we have $x_i^0 - sx_i^* = x_i^*(s - s) = 0$.

Thus, $x' \in \mathbb{R}_+^m \setminus \{0\}$ and the $\{i | x'_i > 0\}$ is a “strict and non-empty subset” of $\{1, \dots, m\}$.

Thus, there exists $\alpha > 0$ such that $\alpha x' \in \Delta^{m-1}$, $B^T(\alpha x') = 0$.

Thus, by part (iii) of lemma 1, $(\alpha x', y^*)$ is an equilibrium for (A, B) .

However, $(\alpha x', y^*)$ is not completely mixed since the $\{i | \alpha x'_i = 0\} \neq \emptyset$, thereby contradicting (A, B) is completely mixed.

Thus, it must be that $x^0 = x^*$.

Similarly, we get $y^0 = y^*$.

Thus, $(x^0, y^0) = (x^*, y^*)$ leading to the desired contradiction and proving the proposition.

Thus, (A, B) has a unique equilibrium. Q.E.D.

The following result leads to interesting consequences for completely mixed bi-matrix games.

Proposition 2: Suppose the $m \times n$ bi-matrix game (A, B) is completely mixed. Suppose (x^*, y^*) is a completely mixed equilibrium for (A, B) . Let $u^* = x^{*T} A y^*$ and $v^* = x^{*T} B y^*$. Then $\{\lambda \in \mathbb{R}^n | A \lambda = 0\} \subset \{t y^* | t \in \mathbb{R}\}$ and $\{\gamma \in \mathbb{R}^m | \gamma^T B = 0\} \subset \{t x^* | t \in \mathbb{R}\}$.

Proof: Suppose the $m \times n$ bi-matrix game (A, B) is completely mixed. Thus, $A \neq 0 \neq B$. Since (x^*, y^*) is a completely mixed equilibrium for (A, B) , by proposition 1, (x^*, y^*) is the unique equilibrium for (A, B) . Since $u^* = x^{*T} A y^*$ and $v^* = x^{*T} B y^*$, (x^*, y^*) is completely mixed implies $A y^* = u^* E^{(m)}$ and $B^T x^* = v^* E^{(n)}$.

Suppose there exists $\lambda \in \mathbb{R}^n \setminus \{t y^* | t \in \mathbb{R}\}$ such that $A \lambda = 0$. Since $y^* \gg 0$ (i.e., all coordinates of y^* are strictly positive) for $\varepsilon > 0$ sufficiently small $y^* + \varepsilon \lambda \gg 0$ and $A(y^* + \varepsilon \lambda) = 0$.

Thus, there exists $\alpha > 0$ such that $\alpha(y^* + \varepsilon \lambda) \in \Delta^{n-1}$.

Let $y^0 = \alpha(y^* + \varepsilon \lambda)$. Then, $A y^0 = \alpha A(y^* + \varepsilon \lambda) = \alpha A y^* + \alpha \varepsilon A \lambda = \alpha A y^* + 0 = \alpha u^* E^{(m)}$, since $A y^* = u^* E^{(m)}$.

Since $x^{*T} B = (x^{*T} B y^*) E^{(n)T} = v^* E^{(n)T}$, it follows from part (iii) of lemma 1 that (x^*, y^0) is an equilibrium for (A, B) .

By proposition 2, we must have $y^0 = y^*$.

$y^0 = \alpha(y^* + \varepsilon \lambda) = y^*$ implies $\lambda = \frac{1-\alpha}{\varepsilon} y^*$, contradicting $\lambda \in \mathbb{R}^n \setminus \{t y^* | t \in \mathbb{R}\}$.

Thus, $\{\lambda \in \mathbb{R}^n | A \lambda = 0\} \subset \{t y^* | t \in \mathbb{R}\}$.

A similar proof works for proving $\{\gamma \in \mathbb{R}^m | \gamma^T B = 0\} \subset \{t x^* | t \in \mathbb{R}\}$. Q.E.D.

An immediate consequence of proposition 2 is the following result which includes a generalization of part (i) of theorem 1 in Oviedo (1996).

Proposition 3: Suppose the $m \times n$ bi-matrix game (A, B) is completely mixed. Then:

(i) Every collection of $n-1$ columns of A are linearly independent and the rank of the matrix A is at least $n - 1$.

(ii) Every collection of $m - 1$ rows of B are linearly independent and the rank of the matrix B is at least $m - 1$.

(iii) $n + 1 \geq m \geq n-1$.

(iv) If rank of the matrix A is $n-1$ “and” rank of the matrix B is $m-1$, then $m = n$.

(v) If for an equilibrium (x, y) of (A, B) it is the case that $V(\text{rowl}(A, B), (x, y)) = 0 = V(\text{columnl}(A, B), (x, y))$, then rank of the matrix A is $n-1$ and the rank of the matrix B is $m-1$.

Proof: Since (A, B) is completely mixed, by proposition 1, it has a unique completely mixed equilibrium (x^*, y^*) . Thus, $x_i^* > 0$ for all $i \in \{1, \dots, m\}$ and $y_j^* > 0$ for all $j \in \{1, \dots, n\}$.

(i) Towards a contradiction suppose a subset of $n-1$ columns of A are linearly dependent.

Without loss of generality suppose the first $n-1$ columns of A are linearly dependent.

Thus, there exists $\lambda = (\lambda_1, \dots, \lambda_{n-1}) \in \mathbb{R}^{n-1} \setminus \{0\}$ such that $\sum_{j=1}^{n-1} A^j \lambda_j = 0$.

Thus, $\sum_{j=1}^{n-1} A^j \lambda_j + 0A^n = 0$.

$\lambda = (\lambda_1, \dots, \lambda_{n-1}) \in \mathbb{R}^{n-1} \setminus \{0\}$ implies $\{j \in \{1, \dots, n-1\} | \lambda_j > 0\} \neq \emptyset$.

From, proposition 2 we know that there exists $t \in \mathbb{R}$, such that $\lambda_j = ty_j^*$ for $j \in \{1, \dots, n-1\}$ and $0 = ty_n^*$.

Since, $y_j^* > 0$ for all $j \in \{1, \dots, n\}$ and $\{j \in \{1, \dots, n-1\} | \lambda_j \neq 0\} \neq \emptyset$, it must be the case that $t \neq 0$.

On the other hand $0 = ty_n^*$ and $y_n^* > 0$ implies, $t = 0$, thereby leading to a contradiction.

Hence every subset of $n - 1$ columns of A must be linearly independent.

This proves (i).

(ii) The proof of (ii) is exactly as in (i) with B^T playing the role of A and x^* playing the role of y^* .

(iii) From (i) it follows that the rank of the matrix A is at least $n-1$.

Since rank of A is equal to the rank of A^T , it must be that A^T has at least $n - 1$ columns and all subsets of $n - 1$ columns of A^T are linearly independent.

Since A is a $m \times n$ matrix, A^T is a $n \times m$ matrix.

Thus, it must be the case that $m \geq n - 1$.

From (ii) it follows that the rank of the matrix B^T is at least $m-1$.

Since rank of B is equal to the rank of B^T , it must be that B has at least $m - 1$ columns and all subsets of $m - 1$ columns of B are linearly independent.

Since B is a $m \times n$ matrix, it must be the case that $n \geq m - 1$.

Thus, $n + 1 \geq m \geq n - 1$.

(iv) Suppose the rank of the matrix A is $n-1$ and rank of the matrix B is $m-1$.

Let (x^*, y^*) be a completely mixed equilibrium for (A, B) . Let $u^* = x^{*T} A y^*$ and $v^* = x^{*T} B y^*$.

We know by the Complementary Slackness Condition that $Ax^* = u^* E^{(m)}$ and $B^T x^* = v^* E^{(n)}$.

Towards a contradiction suppose either $m = n - 1$ or $n = m - 1$.

First suppose $m = n - 1$.

Since rank of the matrix A is $n - 1$, without loss of generality suppose that the first $n - 1$ columns of A are linearly independent. Let A^{-n} denote the $m \times m$ submatrix of A comprising of the first m columns of A and $y^{*(n)}$ be the m -dimensional column vector comprising the first m coordinates of y^* .

Then, given any $\mu \in \{1, -1\}$, there exists a unique $\lambda(\mu) \in \mathbb{R}^m$ such that $A^{-n}\lambda(\mu) = \mu E^{(m)}$. Clearly, $\lambda(\mu) \neq 0$ and $\lambda(-\mu) = -\lambda(\mu)$ for $\mu \in \{1, -1\}$.

If $\{j \in \{1, \dots, m\} \mid \lambda_j(1) > 0\} \neq \emptyset$, then set $\mu = 1$. Otherwise, set $\mu = -1$.

Let $\lambda = \lambda(\mu)$ for the chosen value of μ .

Thus, $\{j \in \{1, \dots, m\} \mid \lambda_j > 0\} \neq \emptyset$

Thus, there exists $\alpha > 0$ such that $y_j^* - \alpha \lambda_j > 0$ for all $j \in \{1, \dots, m\}$.

Choose any such α .

Thus, $A^{-n}(\alpha \lambda) = \alpha \mu E^{(m)}$.

$\{j \in \{1, \dots, m\} \mid \lambda_j > 0\} \neq \emptyset$.

Let $\varepsilon^* = \max\{\varepsilon \mid y_j^* - \alpha \lambda_j \geq 0 \text{ for all } j \in \{1, \dots, m\}\}$.

Thus, $A^{-n}(\varepsilon^* \alpha \lambda) = \alpha \varepsilon^* \mu E^{(m)}$.

Further, $\{j \mid y_j^* - \alpha \varepsilon^* \lambda_j > 0, j \in \{1, \dots, m\}\}$ is a proper subset of $\{1, \dots, m\}$ and $Ay^* - A^{-n}(\varepsilon^* \alpha \lambda) = \alpha \varepsilon^* E^{(m)} = (u^* - \alpha \varepsilon^* \mu) E^{(m)}$.

Let $z^* \in \mathbb{R}^n$ be such that $z_j^* = y_j^* - \alpha \varepsilon^* \lambda_j$ for $j \in \{1, \dots, n-1\}$ and $z_n^* = y_n^*$.

Thus, $z_j^* \geq 0$ for $j \in \{1, \dots, n\}$ and $z_n^* = y_n^* > 0$. Further, $\{j \in \{1, \dots, m\} \mid z_j^* = 0\} \neq \emptyset$.

Since $\sum_{j=1}^n z_j^* > 0$, there exists $\beta > 0$ such that $\beta z^* \in \Delta^{n-1}$.

$A(\beta z^*) = \beta(u^* - \alpha \varepsilon^* \mu) E^{(m)}$.

Since, $B^T x^* = v^* E^{(n)}$, by part (iii) of lemma 1, $(x^*, \beta z^*)$ is an equilibrium for (A, B) .

Since $\{j \mid \beta z_j^* = 0\} \neq \emptyset$, $(x^*, \beta z^*)$ is not completely mixed contradicting (A, B) is completely mixed.

Hence, it must be the case that $m > n - 1$, i.e., $m \in \{n, n + 1\}$.

By a similar argument, it follows that $n \in \{m, m + 1\}$.

Since, $n + 1 \geq m \geq n - 1$, it follows immediately from $m \in \{n, n + 1\}$ and $n \in \{m, m + 1\}$

(v) Let (x, y) be an equilibrium for (A, B) with $V(\text{row}(A, B), (x, y)) = 0 = V(\text{column}(A, B), (x, y))$.

By hypothesis, (x, y) is completely mixed.

From (i) and (ii) it follows that rank of the matrix A is $n - 1$ and rank of the matrix B is $m - 1$.

Thus from (iv) it follows that $m = n$. Q.E.D.

The following theorem collects together the major results we have obtained so far.

Theorem 1: Let (A, B) be a $m \times n$ completely mixed bi-matrix game. Then $n + 1 \geq m \geq n - 1$ with the rank of A being at least $n - 1$, the rank of B being at least $m - 1$, and there exists real numbers u, v such that the linear system $[Ay = uE^{(m)}, B^T x = vE^{(n)}, x \in \Delta^{m-1}, y \in \Delta^{n-1}]$ has a solution. If the rank of the matrix A is $n - 1$ and the rank of the matrix B is $m - 1$, then it must be the case that $m = n$.

Proof: Suppose (A, B) is an $m \times n$ completely mixed bi-matrix game. Then, by proposition 1 we know that it has a unique equilibrium (x^*, y^*) that is completely mixed. By the Complementary Slackness Condition, $Ay^* = (x^{*T}Ay^*)E^{(m)}$, $B^T x^* = (x^{*T}By^*)E^{(n)}$. By (i) and (ii) of proposition 3, we know that the rank of the matrix A is at least $n - 1$ and the rank of the matrix B is at least $m - 1$. By (iv) of proposition 3, we get that if the rank of the matrix A is $n - 1$ and the rank of the matrix B is $m - 1$, then it must be the case that $m = n$. Q.E.D.

5. Completely Mixed TPZS games: An immediate corollary of theorem 1 is the following result for completely mixed TPZS games.

Corollary of theorem 1: Let A be a $m \times n$ TPZS game with $V(A) = 0$. If A is completely mixed, then $m = n$, A has a unique equilibrium and rank of A is $n - 1$.

Proof: We only need to show that rank A is $n - 1$ and $m = n$.

By proposition 1, A has a unique equilibrium (x^*, y^*) that is completely mixed and since $V(A) = 0$, it must be the case that $Ay^* = 0$, $A^T x^* = 0$.

Thus, the columns of A are linearly dependent and the columns of A^T .

Thus, by proposition rank of A is $n - 1$ and rank of A^T is $m - 1$.

Since rank of A is equal to rank of A^T , $n - 1 = m - 1$, so that $m = n$. Q.E.D.

Using the corollary, we can establish the following alternative version of the main result in Kaplansky (1945).

Theorem 2: Let A be a TPZS game with 'n' columns and suppose $V(A) = 0$. Then A is completely mixed if and only if A is a square matrix with the rank of A being equal to $n - 1$, and the linear system $[Ay = 0, y \in \Delta^{n-1}]$ has a solution.

Proof: From corollary of proposition 3, we know that if A is a completely mixed TPZS game with n columns and $V(A) = 0$, then A is a square matrix with the rank of A equal to $n - 1$, and the linear system $[Ay = 0, y \in \Delta^{n-1}]$ has a solution.

Thus, in view of the corollary of proposition 3, all that we need to show for a TPZS game A with 'n' columns satisfying $V(A) = 0$, is that if A is a square matrix of size 'n' (i.e., $n \times n$ matrix), the rank of A is $n - 1$, and the linear system $[Ay = 0, y \in \Delta^{n-1}]$ has a solution, then A is completely mixed.

Let (x^*, y^*) be an equilibrium of the TPZS game A. Thus, $x^* \in \Delta^{m-1}$, $y^* \in \Delta^{n-1}$, $Ay^* \leq (x^{*T}Ay^*)E^{(m)}$, $A^T x^* \geq (x^{*T}Ay^*)E^{(n)} = 0$, since $V(A) = 0$.

$Ay = 0$ implies $x^T Ay = x^{*T} Ay = 0$ for all $x \in \Delta^{n-1}$.

$A^T x^* \geq 0$ implies $x^{*T} Az \geq 0 = x^{*T} Ay$ for all $z \in \Delta^{n-1}$.

Since $y \in \Delta^{n-1}$, we get (x^*, y) is also an equilibrium of the TPZS A.

Hence, without loss of generality suppose, $y = y^*$.

If $\{j | y_j^* = 0\} \neq \emptyset$, then $\sum_{\{j | y_j^* > 0\}} A^j y_j^* = 0$, with the cardinality of $\{j | y_j^* > 0\}$ being at least 1 and at most $n-1$, contradicts our assumption that the rank of A is $n-1$.

Hence $\{j | y_j^* = 0\} = \emptyset$.

Thus, by the Complementary Slackness Condition, it must be the case that $x^{*T} A = 0$.

If $\{i | x_i^* = 0\} \neq \emptyset$, then $\sum_{\{i | x_i^* > 0\}} A_i x_i^* = 0$, with the cardinality of $\{i | x_i^* > 0\}$ being at least 1 and at most $n-1$, contradicting our assumption that the rank of A is $n-1$.

Thus, it must be the case that $\{i | x_i^* = 0\} = \emptyset$.

Suppose, (x^0, y^0) is an equilibrium for A. Then (x^0, y^*) and (x^*, y^0) are also equilibria of A.

Thus, $\{j | y_j^* = 0\} = \emptyset$ implies by the Complementary Slackness Condition that $x^{0T} A = 0$ and $\{i | x_i^* = 0\} = \emptyset$ implies by the Complementary Slackness Condition $Ay^0 = 0$.

Since rank of A is $n-1$ it follows by an argument similar to that for x^* and y^* that $\{i | x_i^0 = 0\} = \emptyset$ and $\{j | y_j^0 = 0\} = \emptyset$.

Thus, (x^0, y^0) is a completely mixed equilibrium, there by leading to the conclusion that A is a completely mixed TPZS game. Q.E.D.

Note: This is a revised version of an earlier paper entitled “Linear Systems Equivalent to Two-Person Zero Sum Games”. This version generalizes and extends several results in the earlier version.

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